RL from Human Feedback as Game-Solving

Gokul Swamy

what happens to online / offline PFT as a result?

2. What is a more robust criterion for preference aggregation and how can we efficiently optimize it?

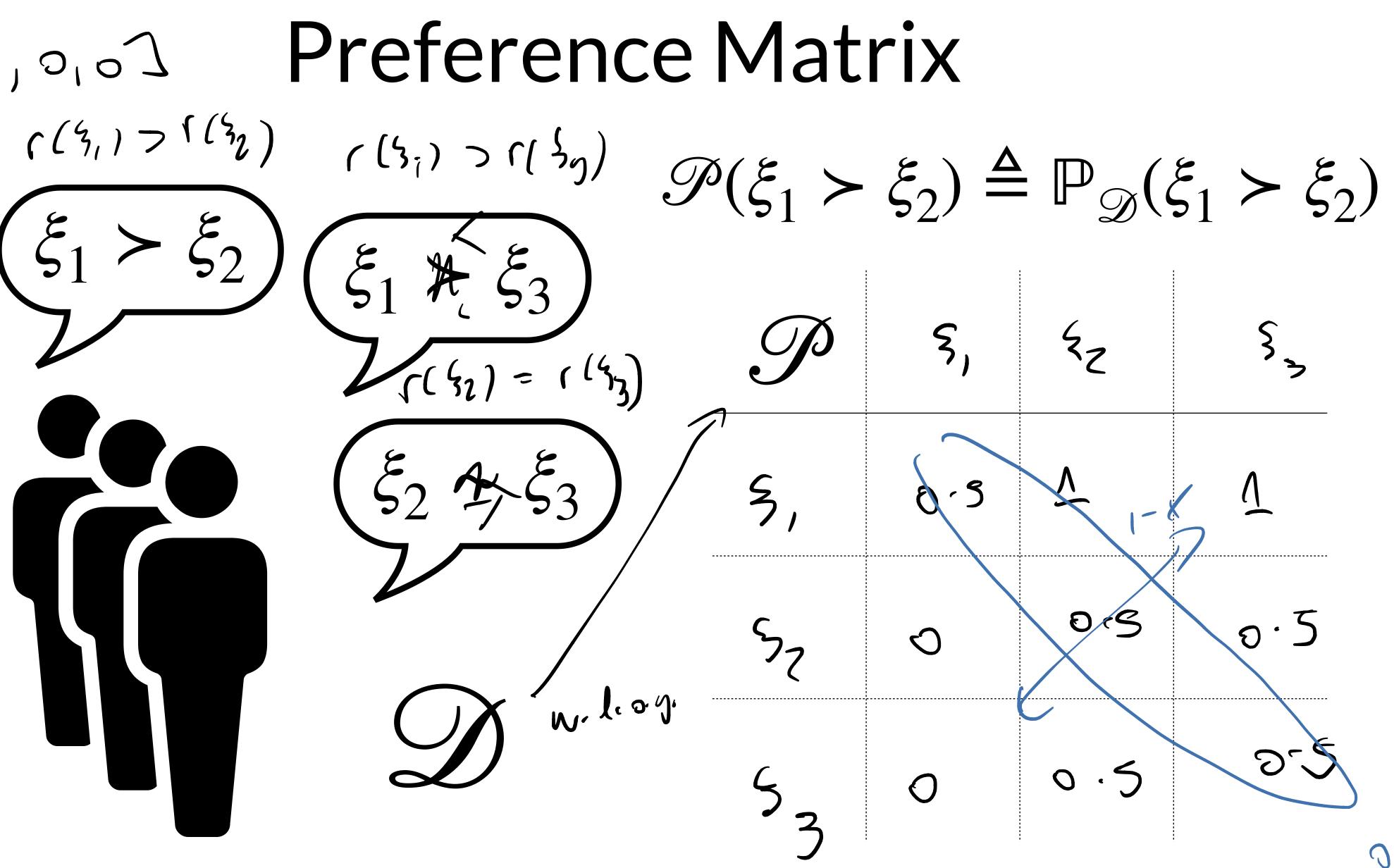
1. When is the Bradley-Terry assumption inaccurate and

1. When is the Bradley-Terry assumption inaccurate and what happens to online / offline PFT as a result?

(aggregate) preferences, leading to mode collapse in RLHF.

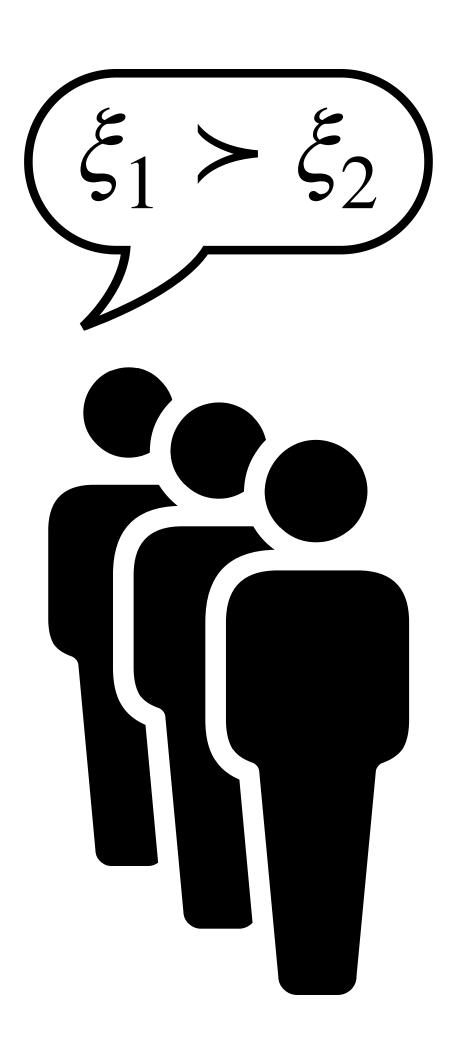
- 2. What is a more robust criterion for preference aggregation and how can we efficiently optimize it?
- **A**: BT is violated when when a reward function can't explain

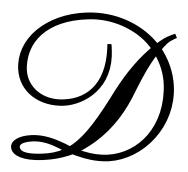
r=CI, 0, 03 Preference Matrix ξ_2 ξ1 ξ3 A r(32) = r(33) ξ_2

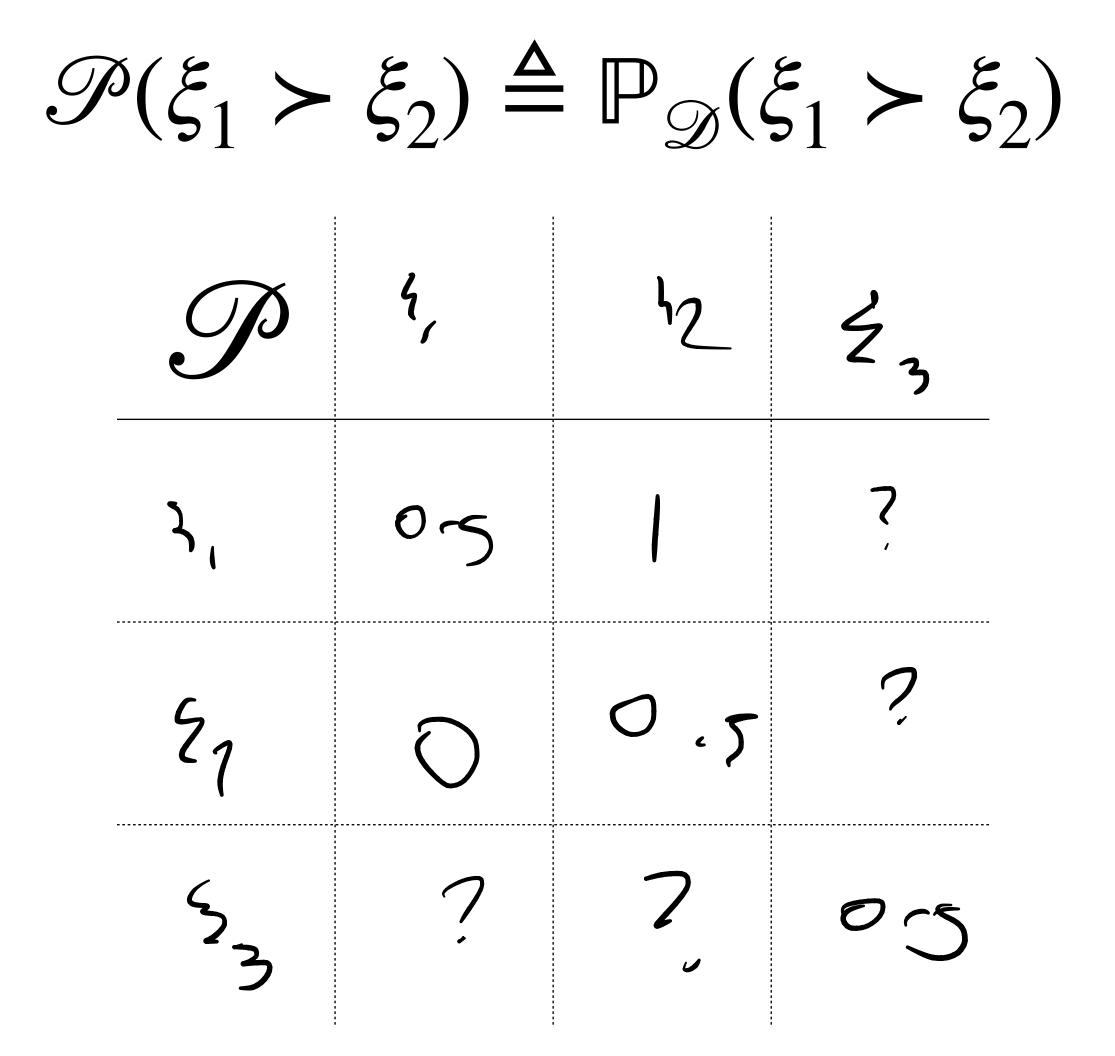




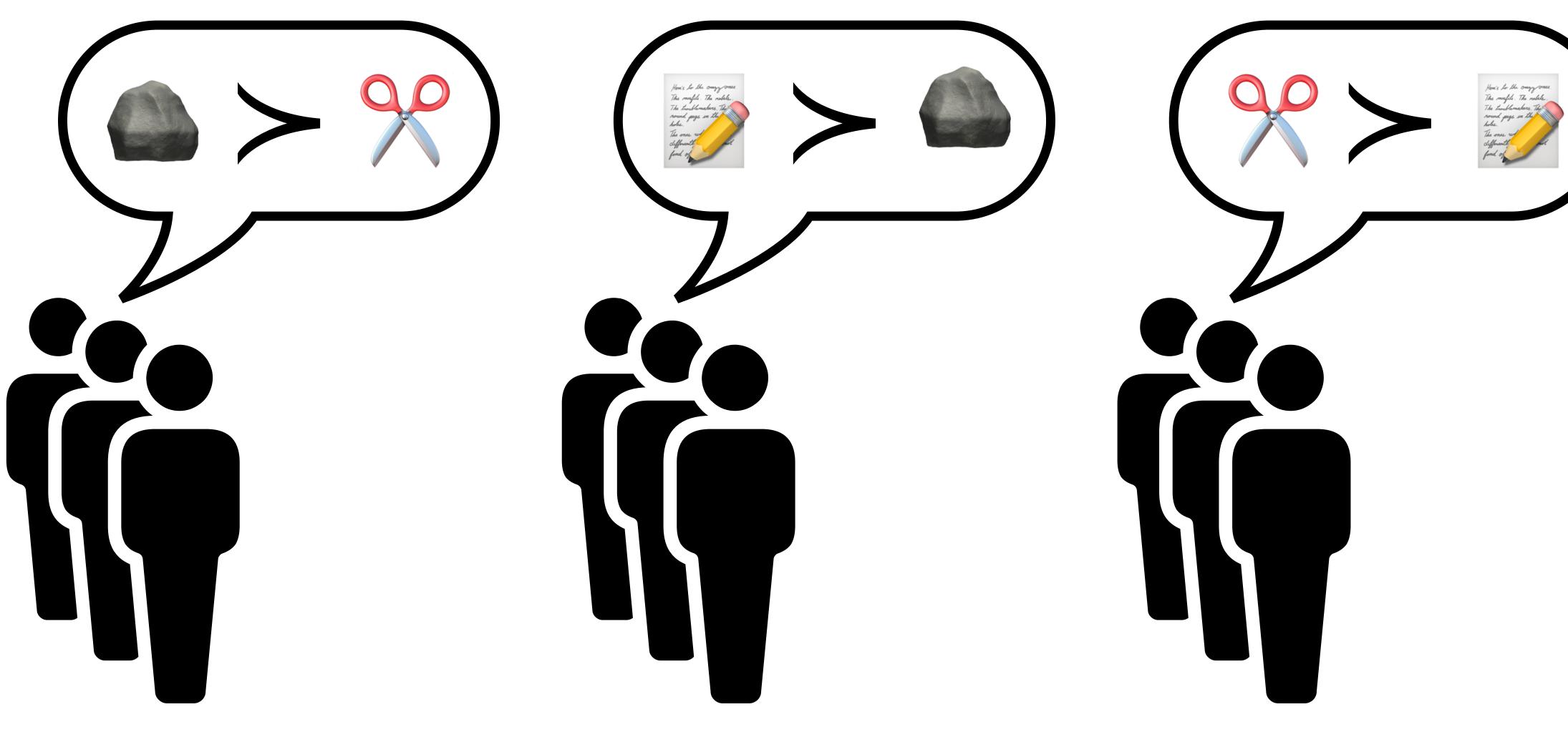
(Partial) Preference Matrix





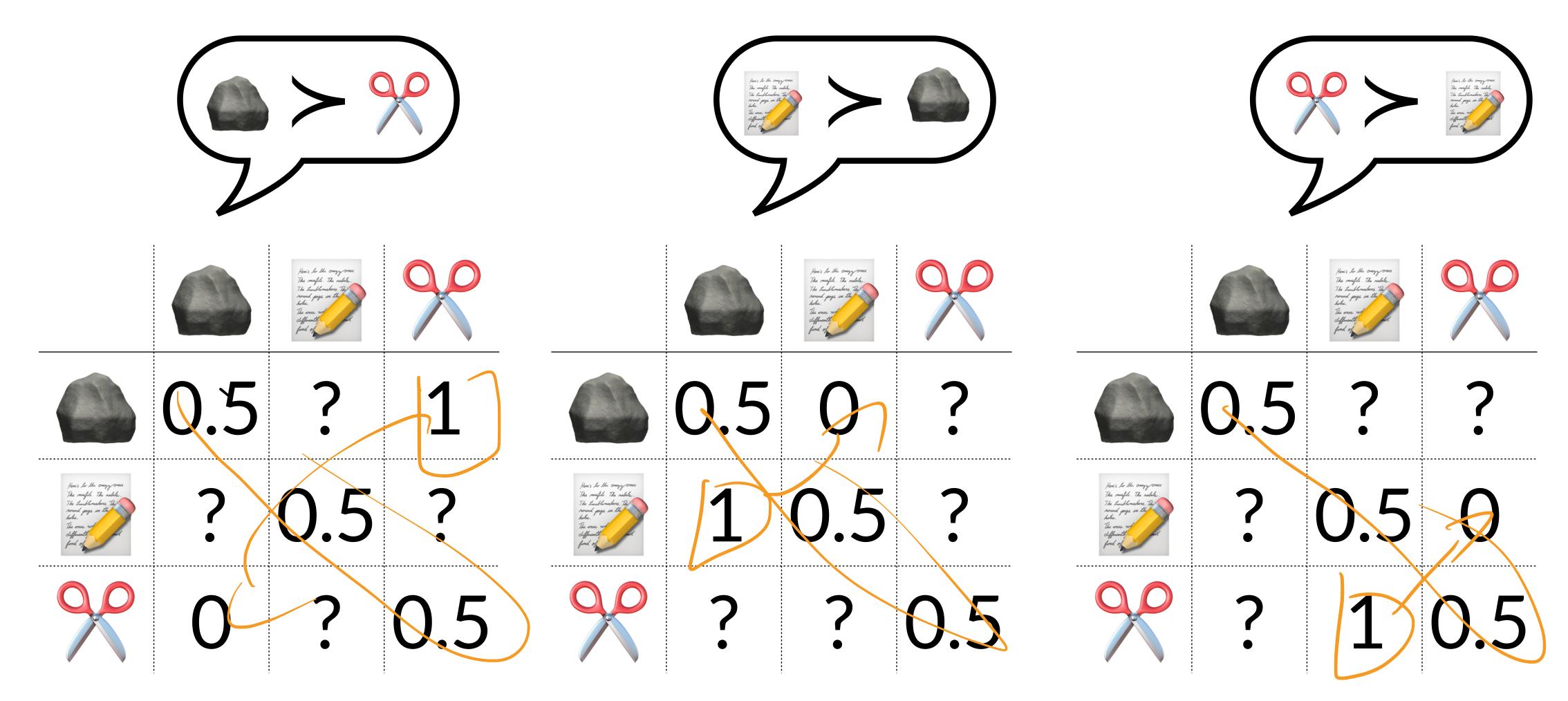


Intransitivity from Preference Aggregation

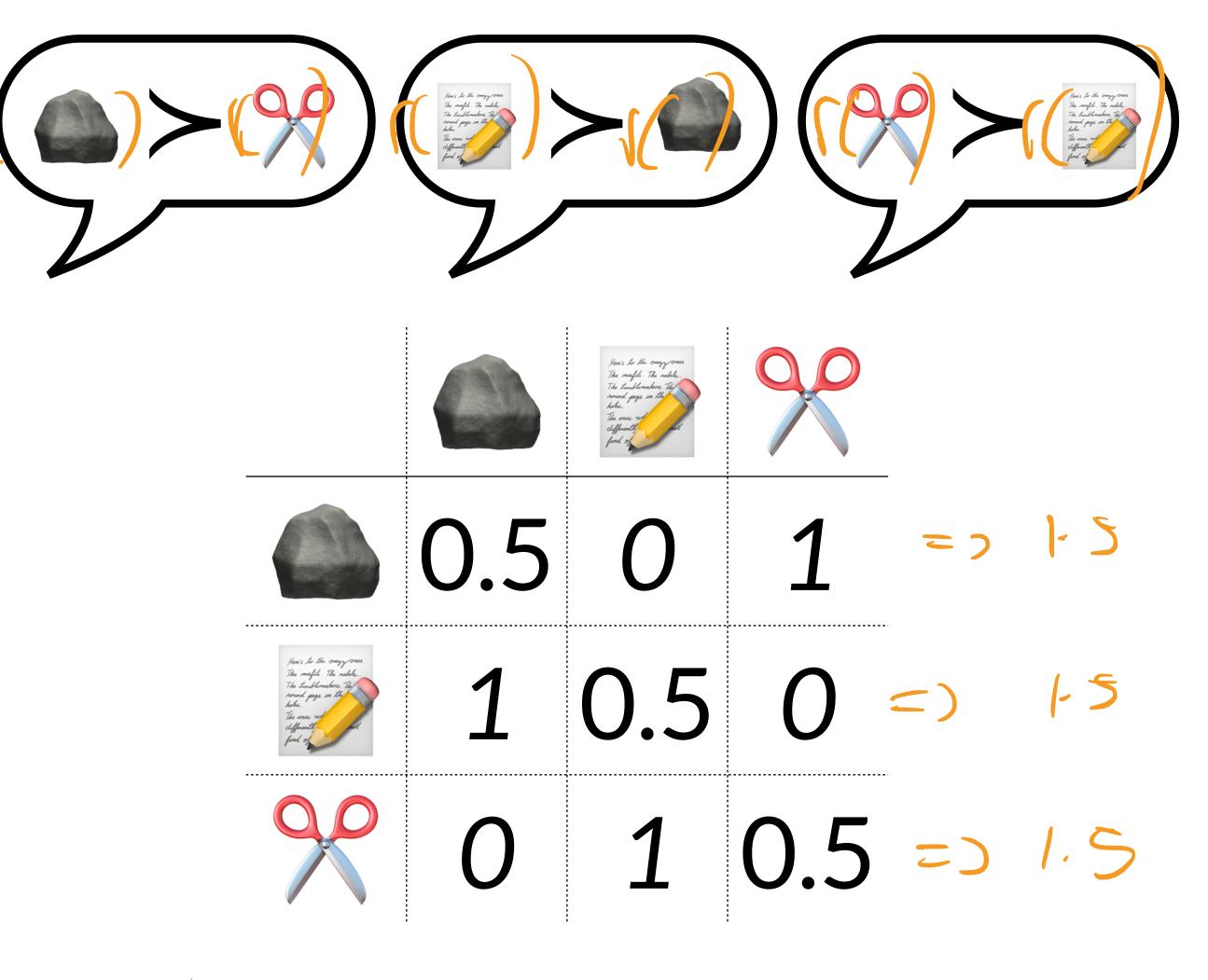




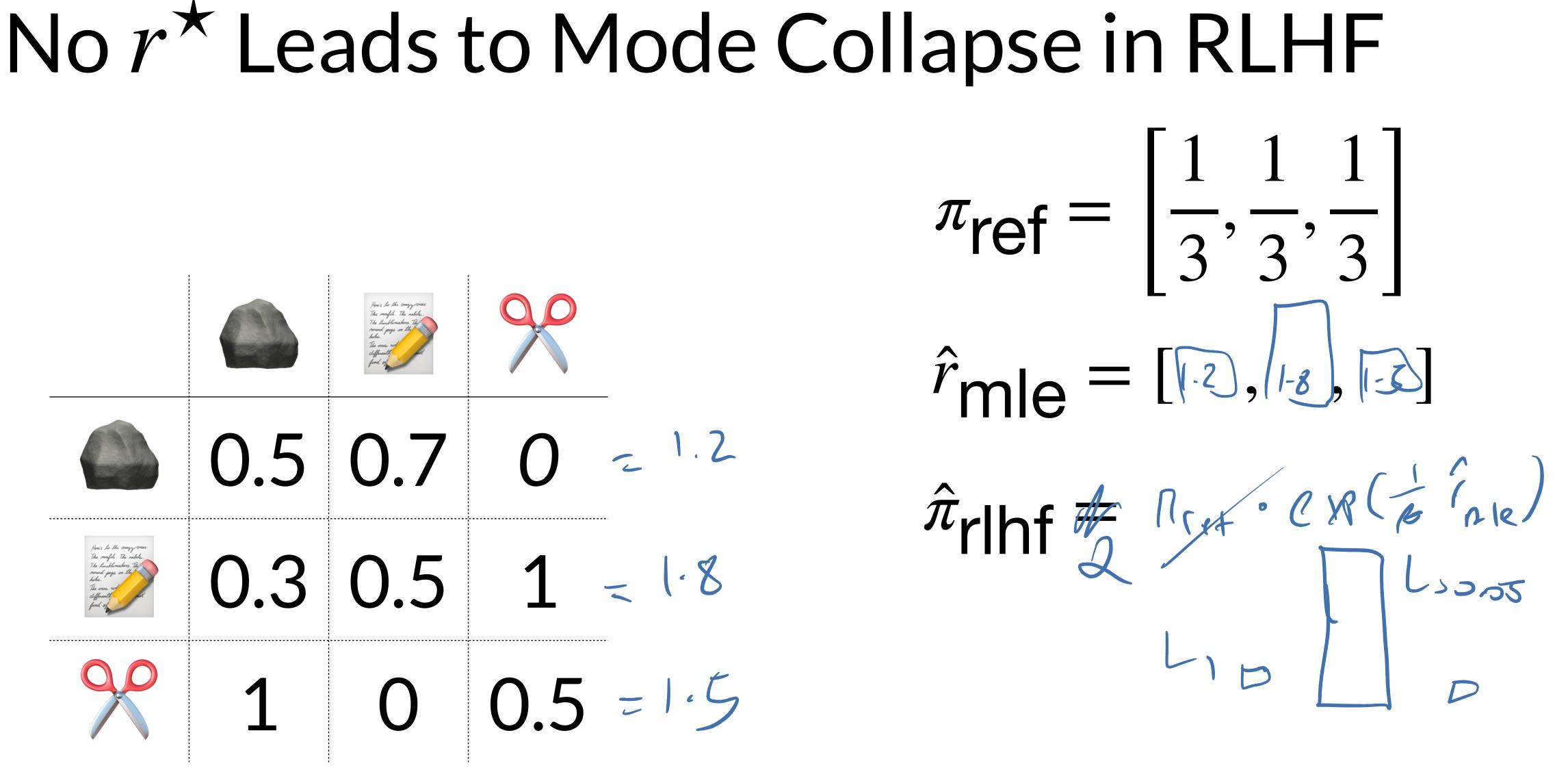
Intransitivity from Preference Aggregation



Intransitivity from Preference Aggregation

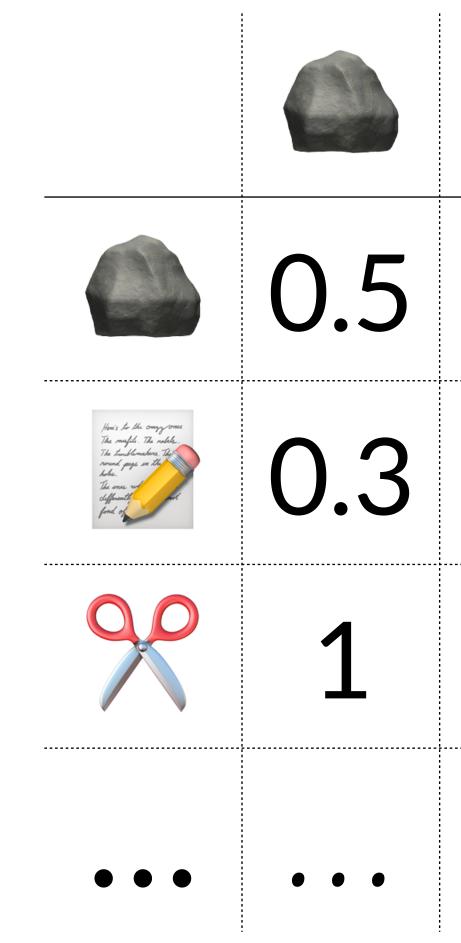


There is no r^* that explains these preferences!

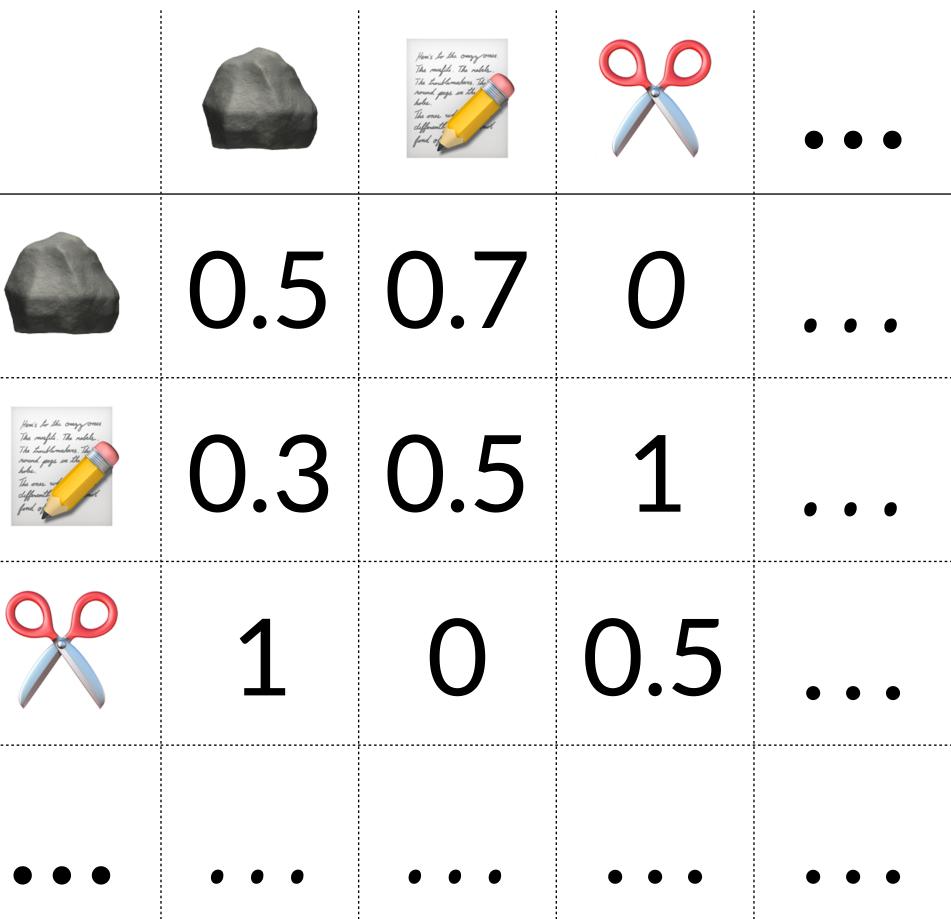




No r* Leads to Mode Collapse in RLHF



This problem gets worse with a larger $\Xi!$



1. When is the Bradley-Terry assumption inaccurate and what happens to online / offline PFT as a result?

(aggregate) preferences, leading to mode collapse in RLHF.

2. What is a more robust criterion for preference aggregation and how can we efficiently optimize it?

A: The minimax winner doesn't assume transitivity of

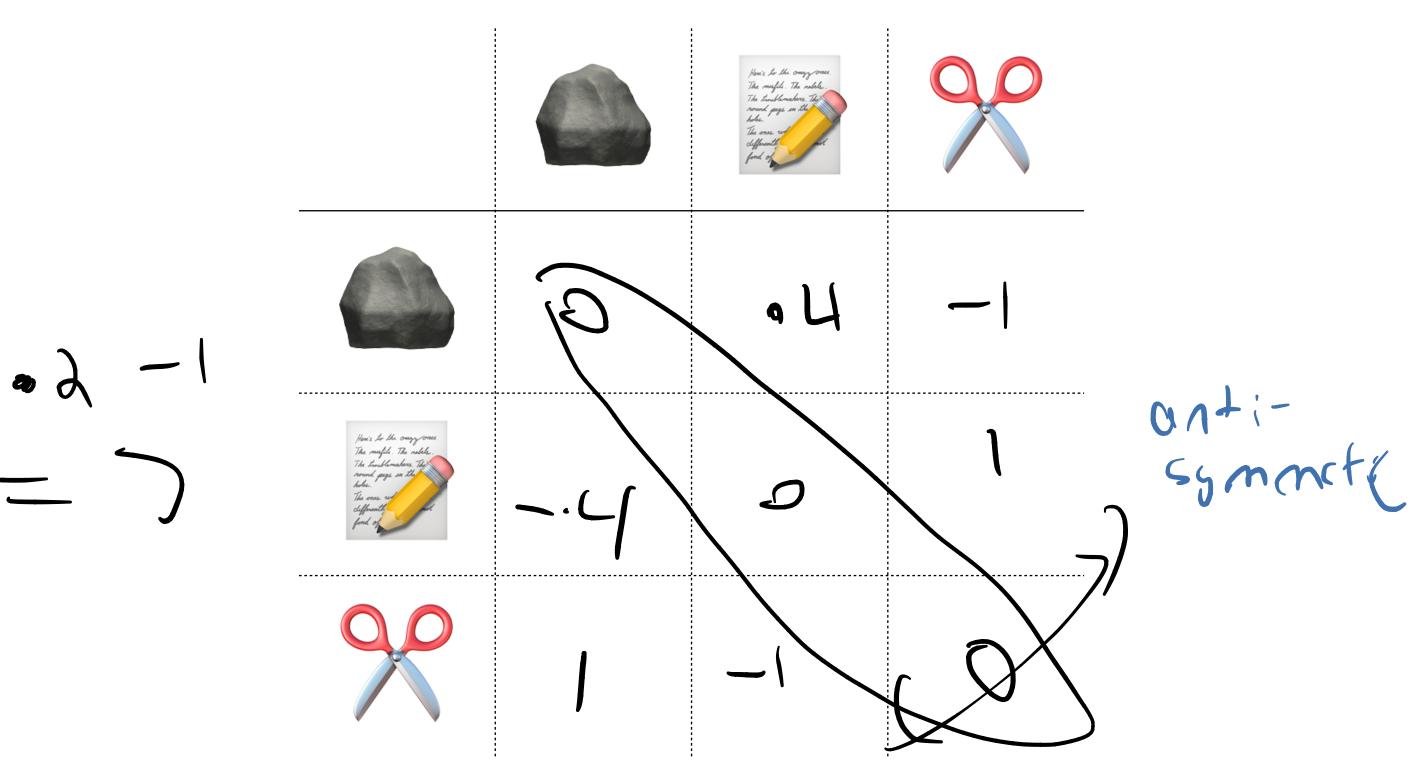
A: BT is violated when when a reward function can't explain

preferences. We can use a self-play algorithm to compute it.

Beyond Bradley-Terry in RLHF



$2 \cdot \mathscr{P} - 1$ defines a symmetric 2p0s game!



Von Neumann / Minimax Winners

 $\pi_1^{\star}, \pi_2^{\star} = \arg\max_{\pi_1 \in \Pi} \arg\min_{\pi_2 \in \Pi} \mathbb{E}_{\xi_1 \sim \pi_1, \xi_2 \sim \pi_2} [2\mathscr{P}(\xi_1 \succ \xi_2) - 1]$

$\max_{n} (n)^{T} (2p_{1}) P_{2}^{*} Z_{n}^{T} (2p_{1}) \Lambda = 0 \quad 2^{T} = 0.5$

- 2. Preferred to any other policy w.p. 1/2.
- 3. No assumptions on underlying shared r^* required!

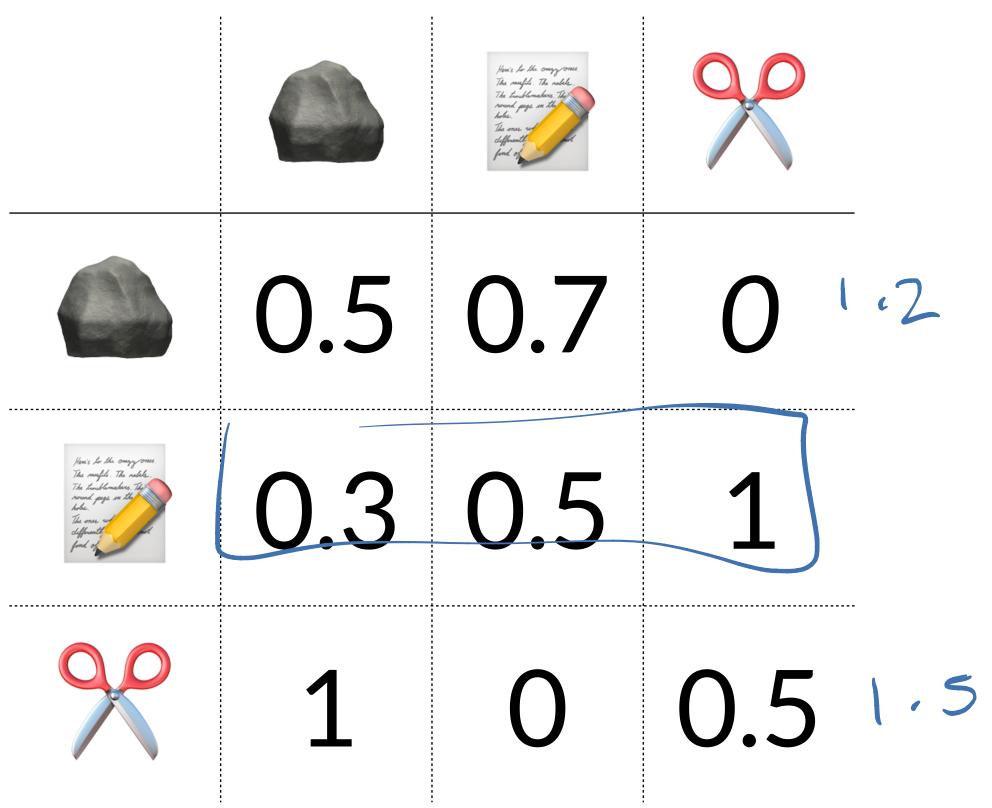
Lo robust against the morst ease comparator policy

1. Pick a policy that is robust against worst case comparator.



Von Neumann / Minimax Winners

$\pi_1^{\star}, \pi_2^{\star} = \arg\max_{\pi_1 \in \Pi} \arg\min_{\pi_2 \in \Pi} \mathbb{E}_{\xi_1 \sim \pi_1, \xi_2 \sim \pi_2} [2\mathscr{P}(\xi_1 \succ \xi_2) - 1]$





_	

If *Computing* Minimax Winners

 $\pi_1^{\star}, \pi_2^{\star} = \arg\max\arg\min$ $\pi_1 \in \Pi \qquad \pi_2 \in \Pi$ We can define a sequence of losses for each NR player: $\ell_t^1(\pi) = \mathbb{E}_{\xi \sim \pi, \xi' \sim \pi_t^t} [2\mathcal{P}(\xi \succ \xi') - 1] \quad \ell_t^2(\pi) = \mathbb{E}_{\xi \sim \pi_t^1, \xi' \sim \pi} [-(2\mathcal{P}(\xi \succ \xi') - 1)]$ Without loss of generality, assume that $\pi_1^0 = \pi_2^0$. Then, $\mathscr{C}_0^1(\pi) = \mathbb{E}_{\xi \sim \pi, \xi' \sim \pi_2^0}[2\mathscr{P}(\xi \succ \xi') - 1]$ $= \mathbb{E}_{\xi \sim \pi, \xi' \sim \pi_1^0} [2\mathscr{P}(\xi \succ \xi') - 1]$ $= \mathbb{E}_{\xi \sim \pi_1^0, \xi' \sim \pi} \left[- \left(2 \mathscr{P}(\xi \succ \xi') - 1 \right) \right] = \ell_0^2(\pi)$

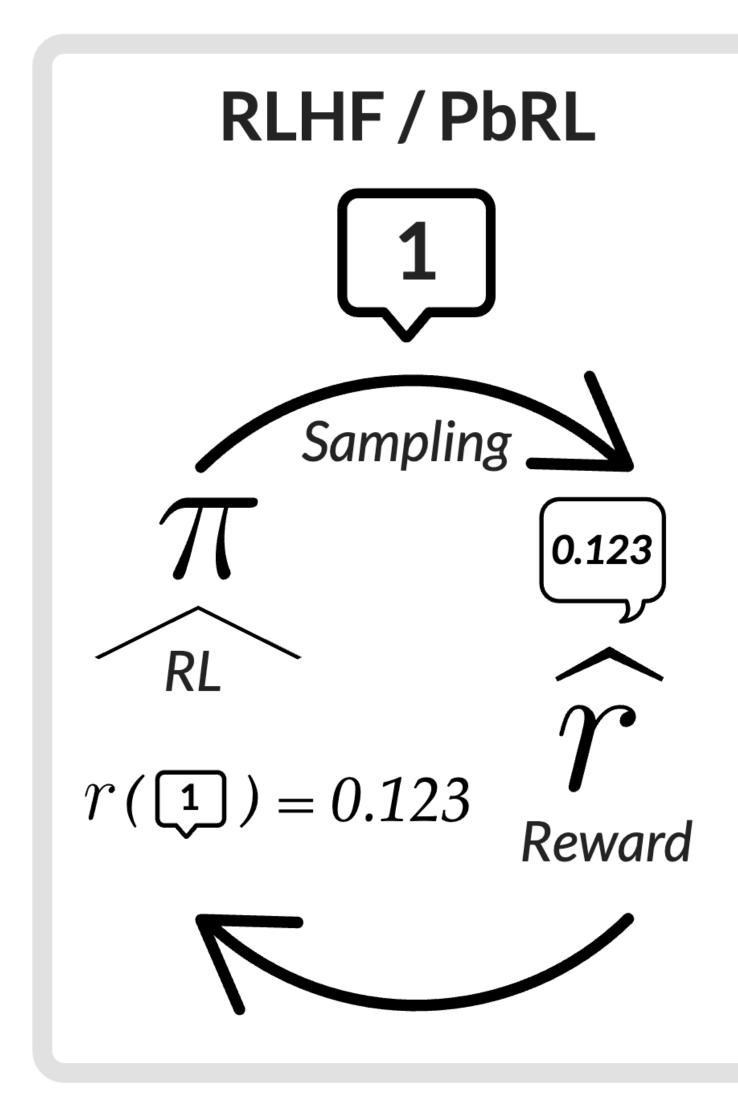
 $\Rightarrow \forall t \in [T], \pi_1^t = \pi_2^t$

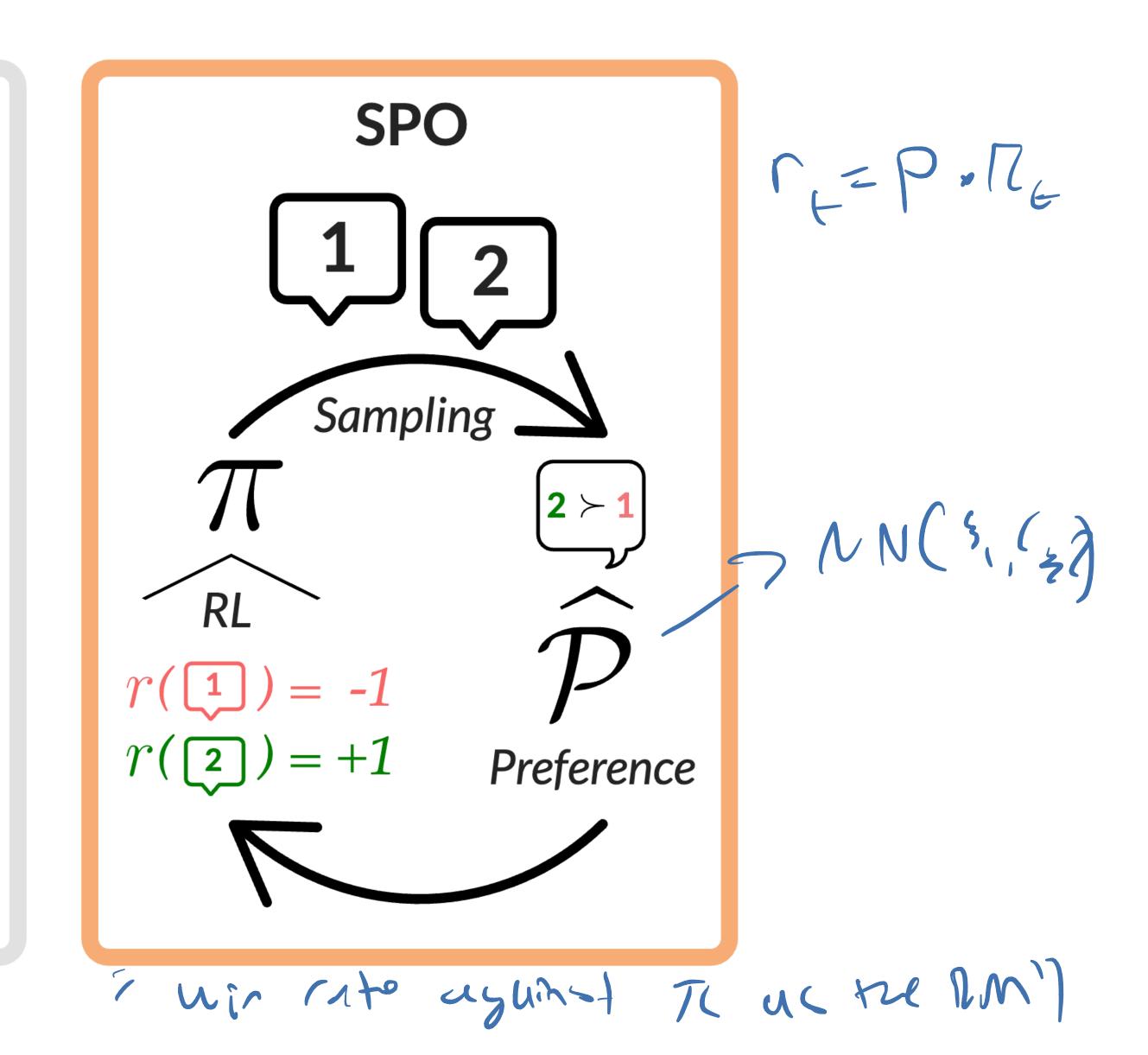
Can just do self-play, no adversarial training required!

$$\mathbb{E}_{\xi_1 \sim \pi_1, \xi_2 \sim \pi_2} [2\mathcal{P}(\xi_1 > \xi_2) - 1]$$



SPO: Self-Play Preference Optimization





Offline Dataset P(r > c) SFT RLHF SPO P(r > c) RLHF ISPO DSPO 0.02 0.02 RLHF 0.5 0.21 0.24 0.5 RLHF 0.98 0.5 0.25 ISPO 0.79 0.5 0.61 0.98 0.75 0.5 DSPO 0.76 0.39 0.5 SPO

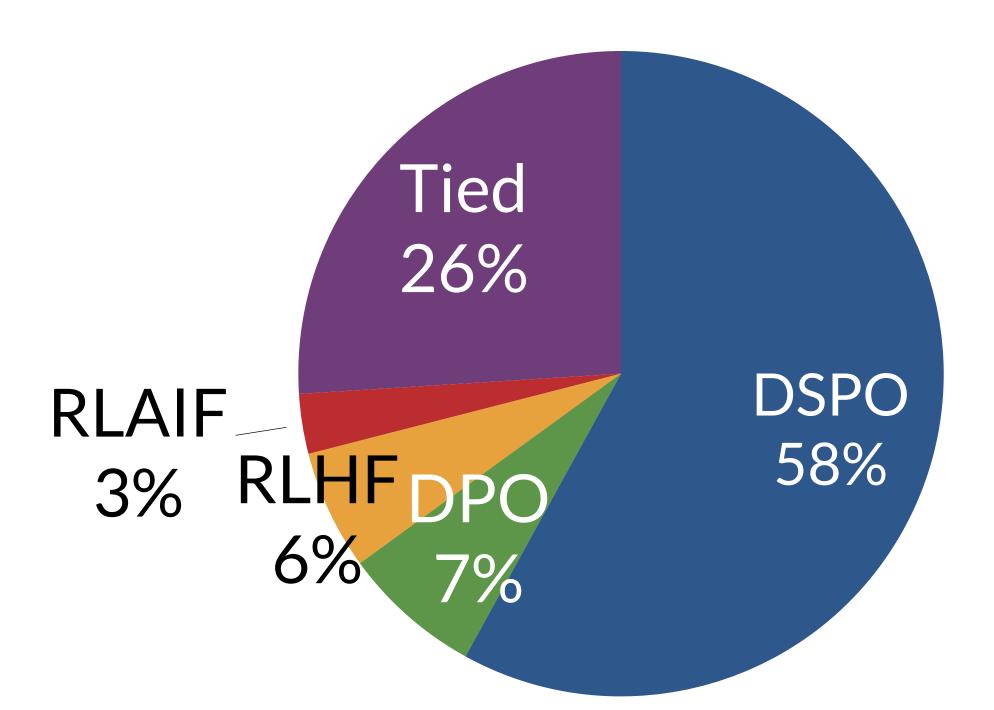
[Munos+'23]

Instantiations of SPO on Large Language Models

[Calandriello+'24]



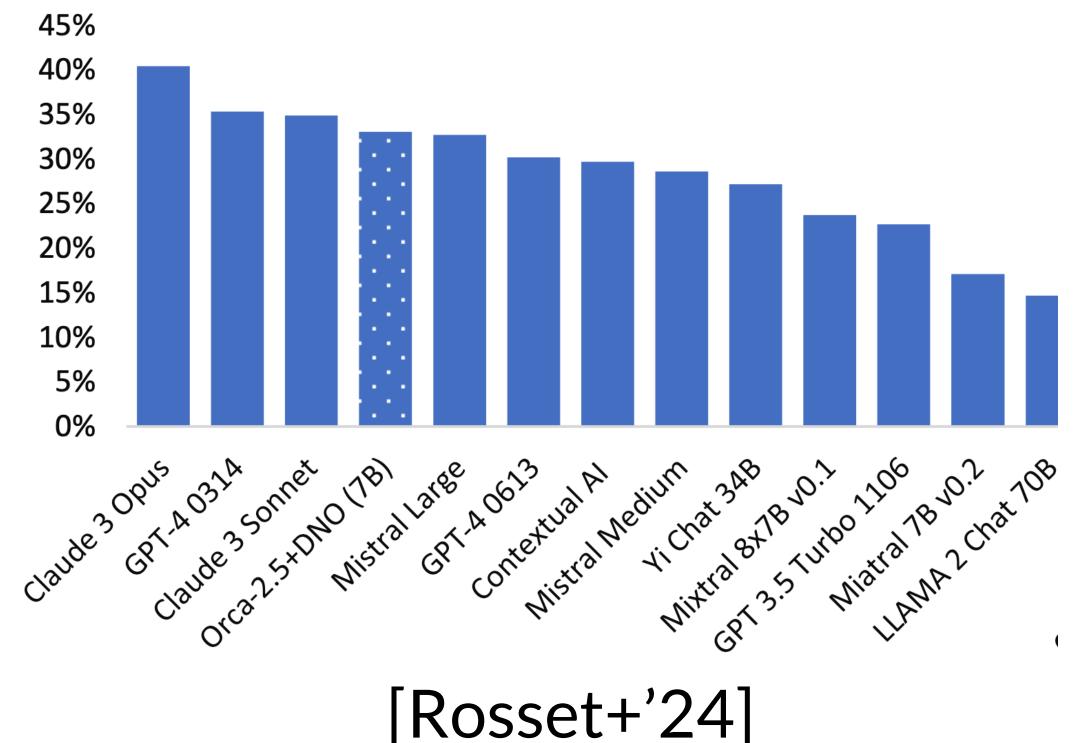
Instantiations of SPO on Large Language Models



[Guo+'24]

Online Oracle

LC Win Rate



1. When is the Bradley-Terry assumption inaccurate and what happens to online / offline PFT as a result?

(aggregate) preferences, leading to mode collapse in RLHF.

2. What is a more robust criterion for preference aggregation and how can we efficiently optimize it?

A: The minimax winner doesn't assume transitivity of

A: BT is violated when when a reward function can't explain

preferences. We can use a self-play algorithm to compute it.